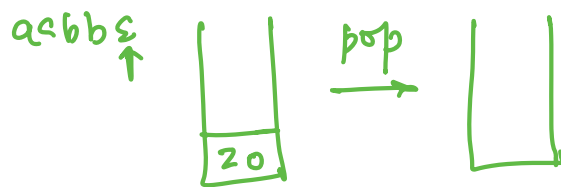
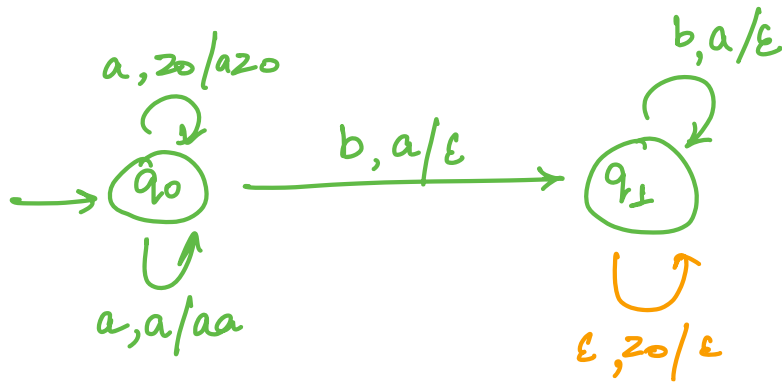
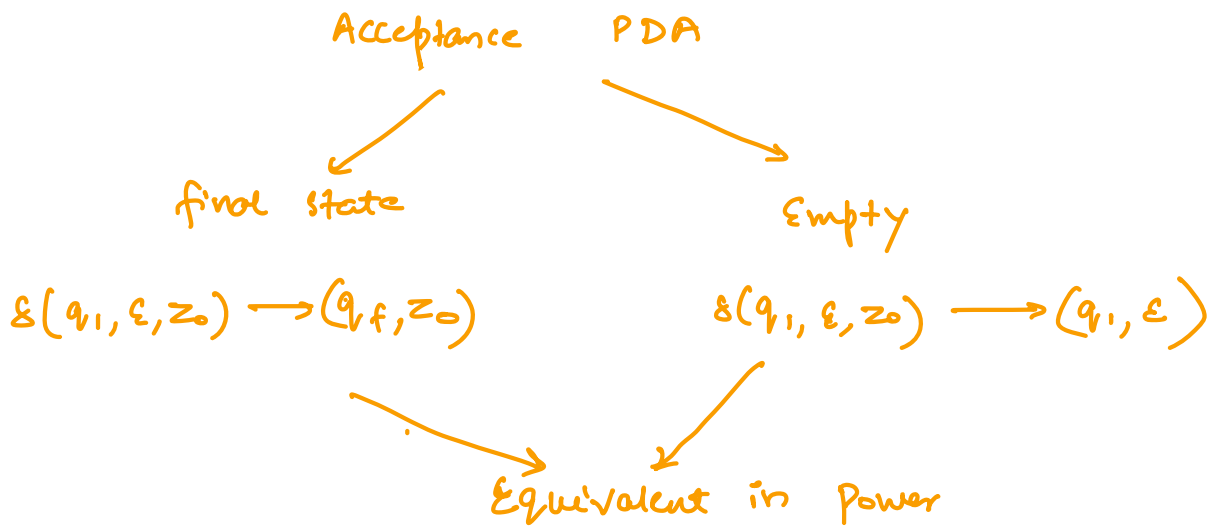


Acceptance by Empty Stack



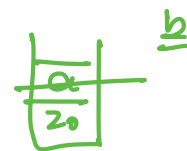
Ⓐ

Ⓑ | $\delta(q_1, \epsilon, z_0) \rightarrow (q_1, \epsilon)$



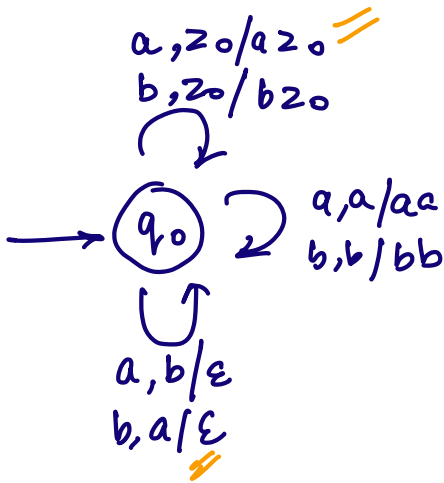
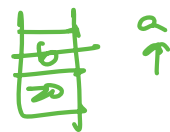
Eq: $w \mid n_a(w) = n_b(w)$

Logic: first time : push a, push b
 top a, input b pop



top b, input a

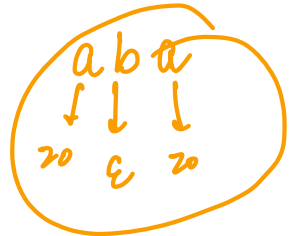
pop



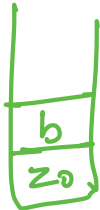
$\epsilon, z_0 / z_0$



why not a new state for pop?



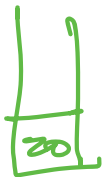
baab ϵ



bbaa ϵ



baab ϵ



bbaa ϵ



baab ϵ



bbaa ϵ



baab ϵ



bbaa ϵ



baabε
↑



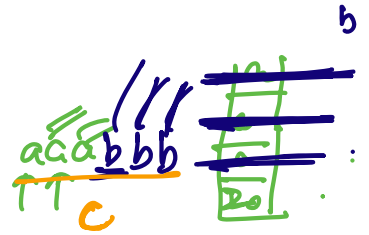
bbaaε
↑



Eg: $a^n b^n c^m \mid n, m \geq 1$

a: push
b: pop

at last 1 c → final state



(Push a on empty stack)

a, z0/z0



b, a/ε (pop)

b, a/ε (pop)



c, z0/z0
first c

c, z0/z0



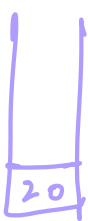
more no of c's

Push a's

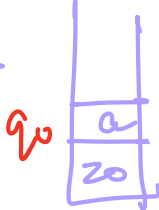
abc

aabc

aaabc



→



→



aaabc

aaabc



dead state
no transition available X

$\delta(q_1, c a) ?$
X

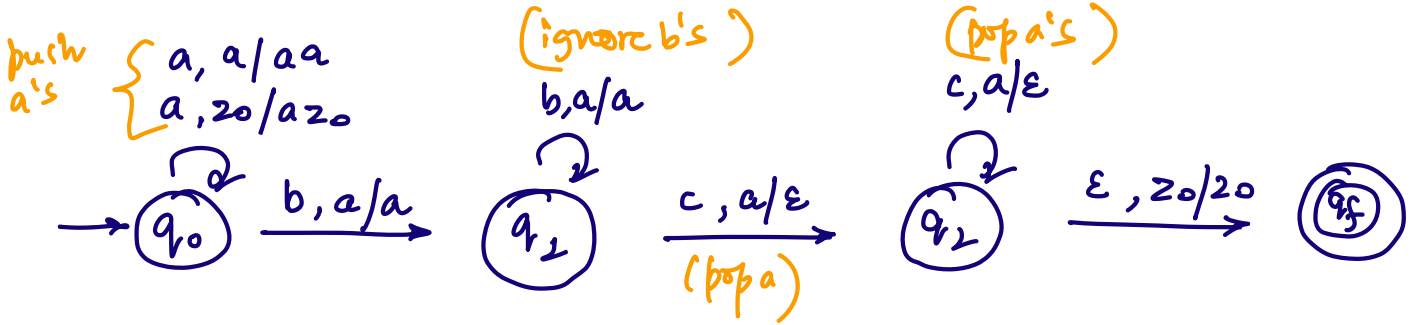
DFA:



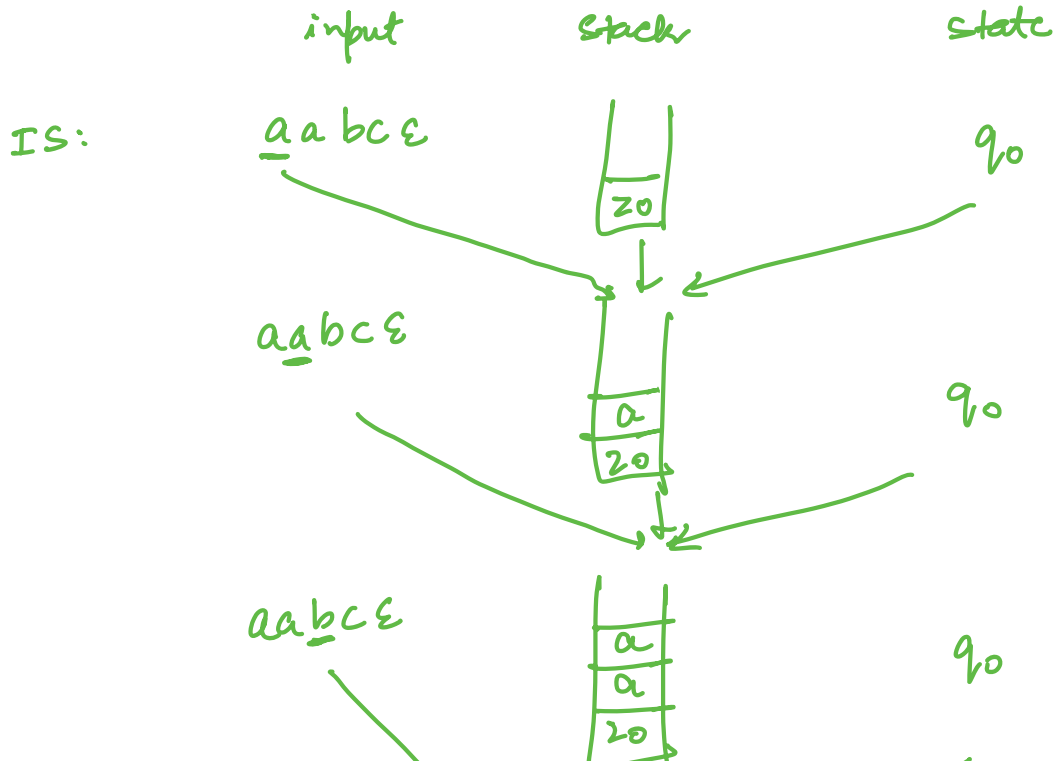
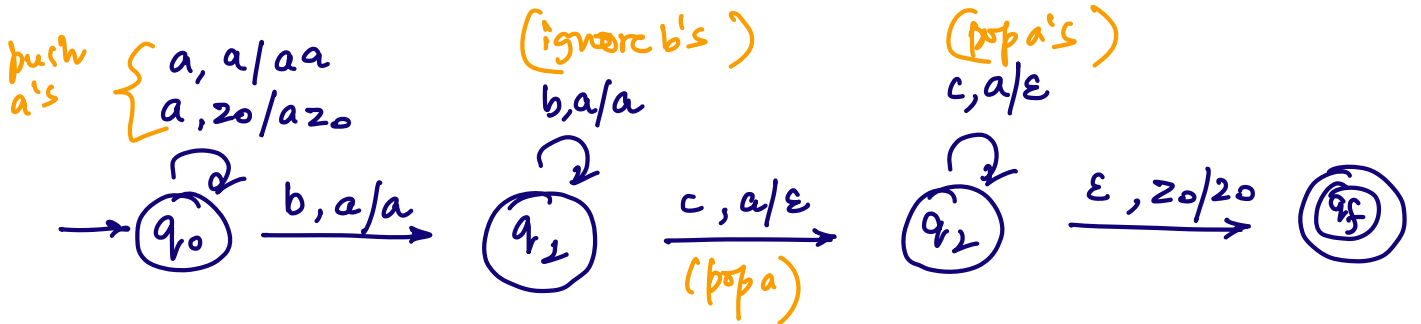
eg: $a^n b^m c^n \mid n, m \geq 1$

alphabet, stack top / new top

a: push
b: ignore
c: pop



aa b c ϵ



aabcε



q₁

aabcε



q₂

$\delta(q_2, \epsilon, a) ? \times$

Dead State \rightarrow String not accepted

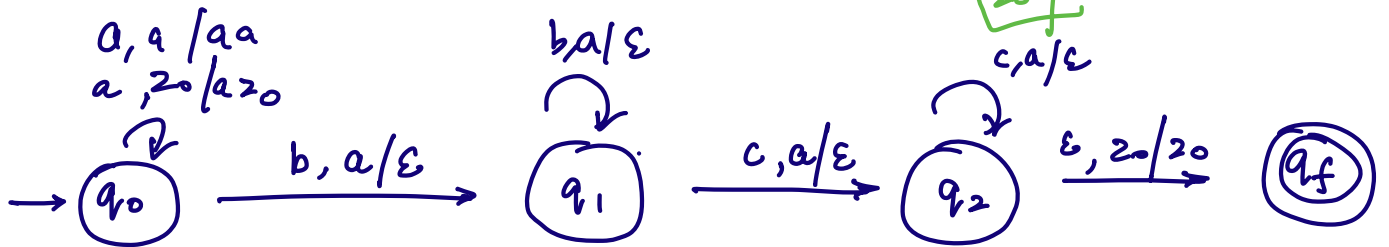
eg: $a^{m+n} b^m c^n \mid n, m \geq 1$

a: push
b: pop
c: pop

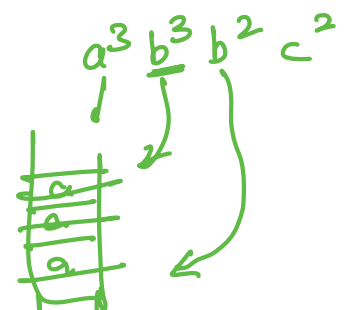
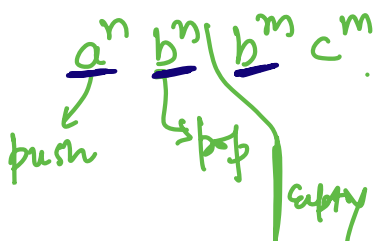
aaaa bb ccc

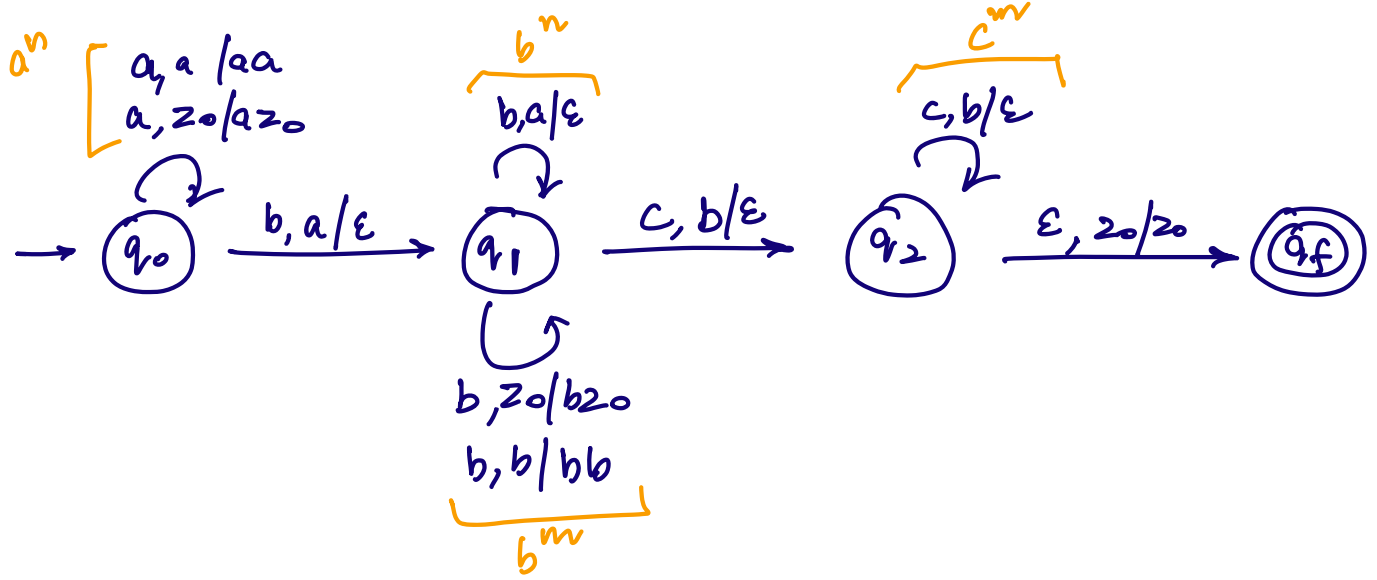


Stack at state q₂

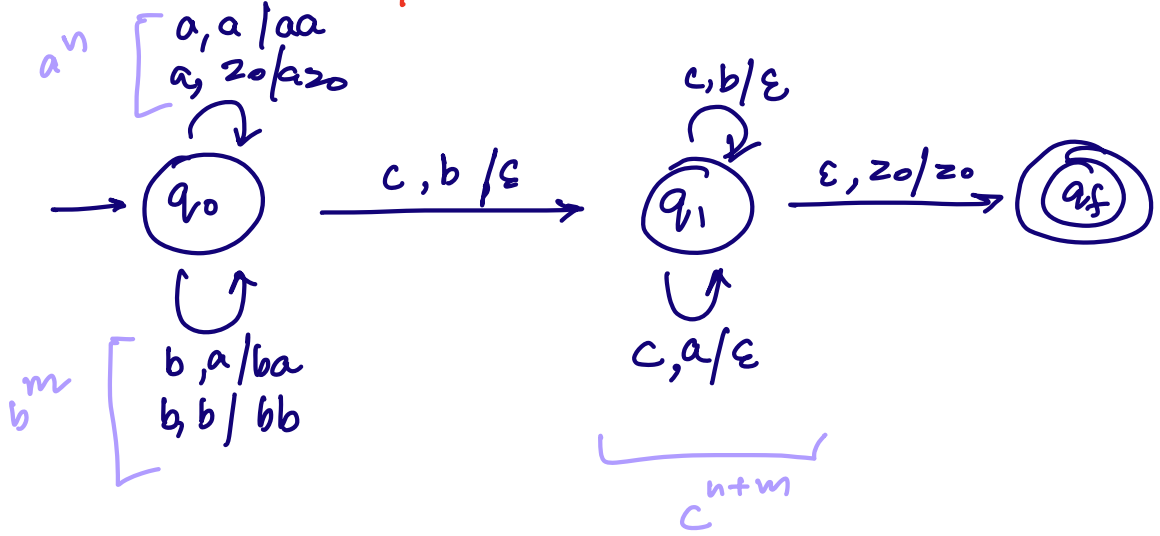


eg: $a^n b^{n+m} c^m \mid n, m \geq 1$

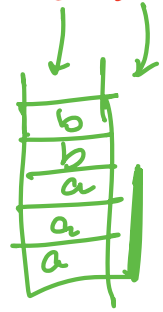




Eq: $a^n b^m c^{n+m} \mid n, m \geq 1$

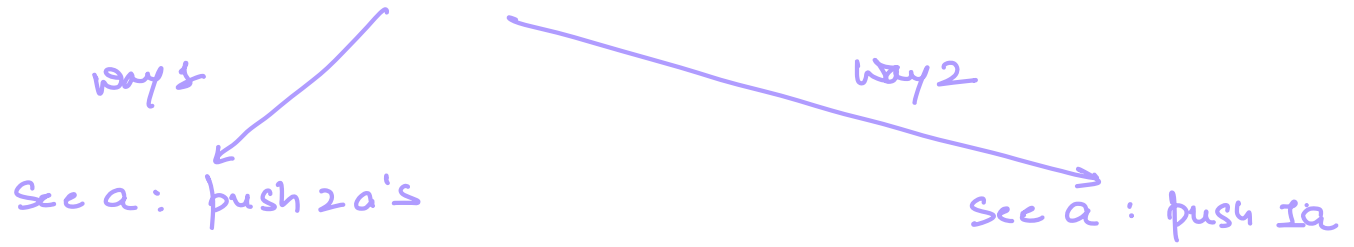


Eq: $a^n b^m c^n d^m \mid n, m \geq 1$



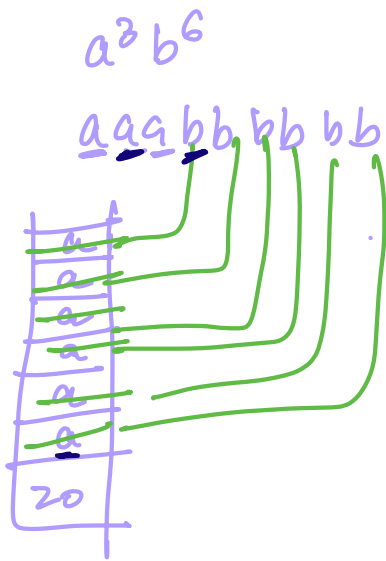
Single stack PDA not possible

Eq: $a^n b^{2n} \mid n \geq 1$

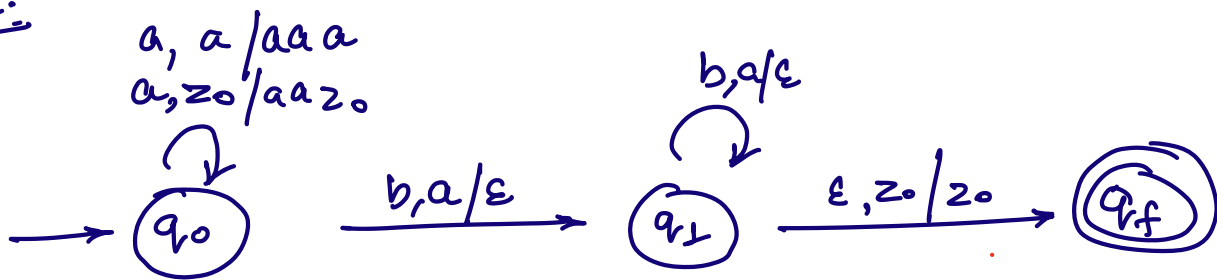


See b : pop 1 a

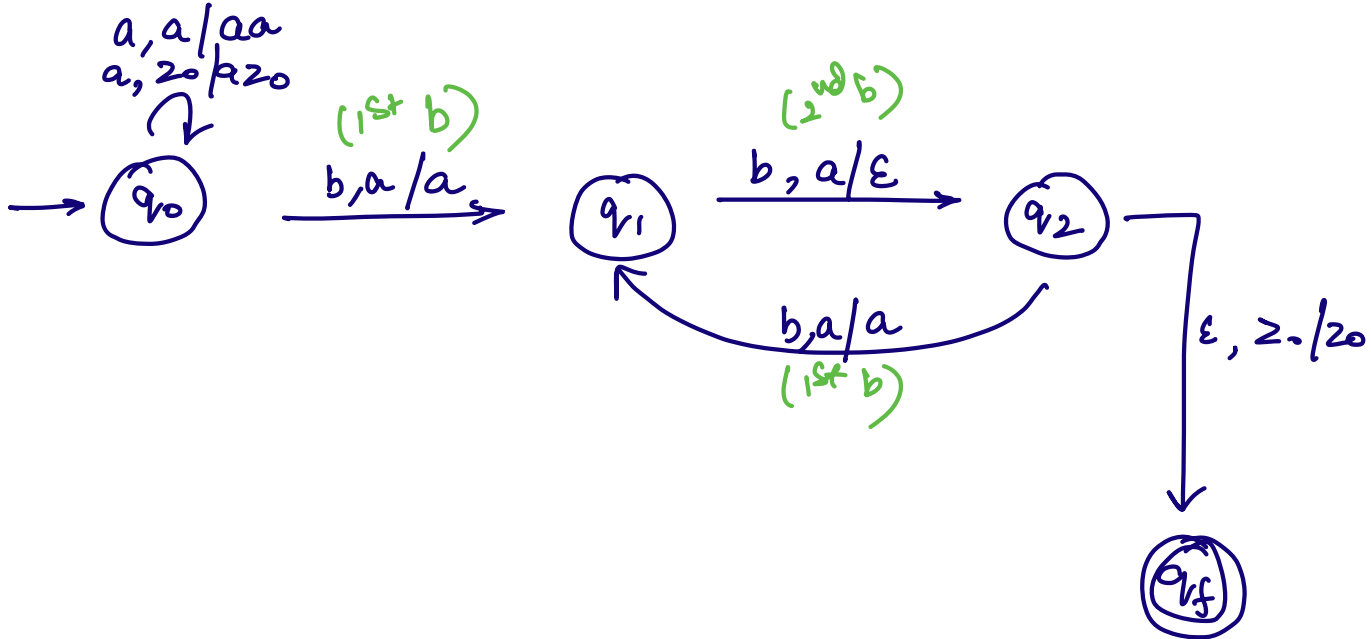
See 2 b's : pop 1 a



Way 1:



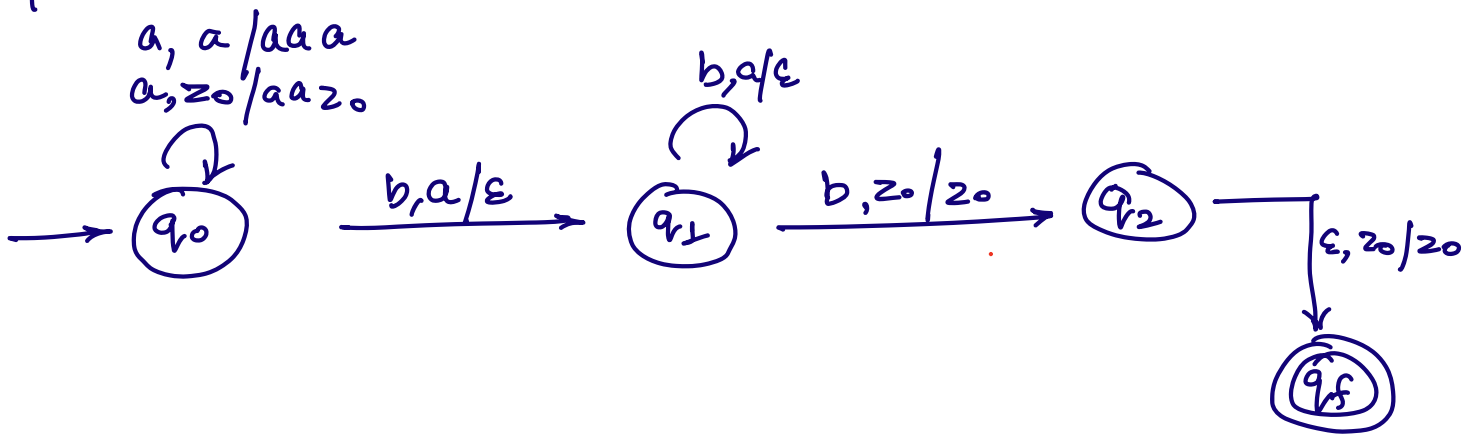
Way 2:



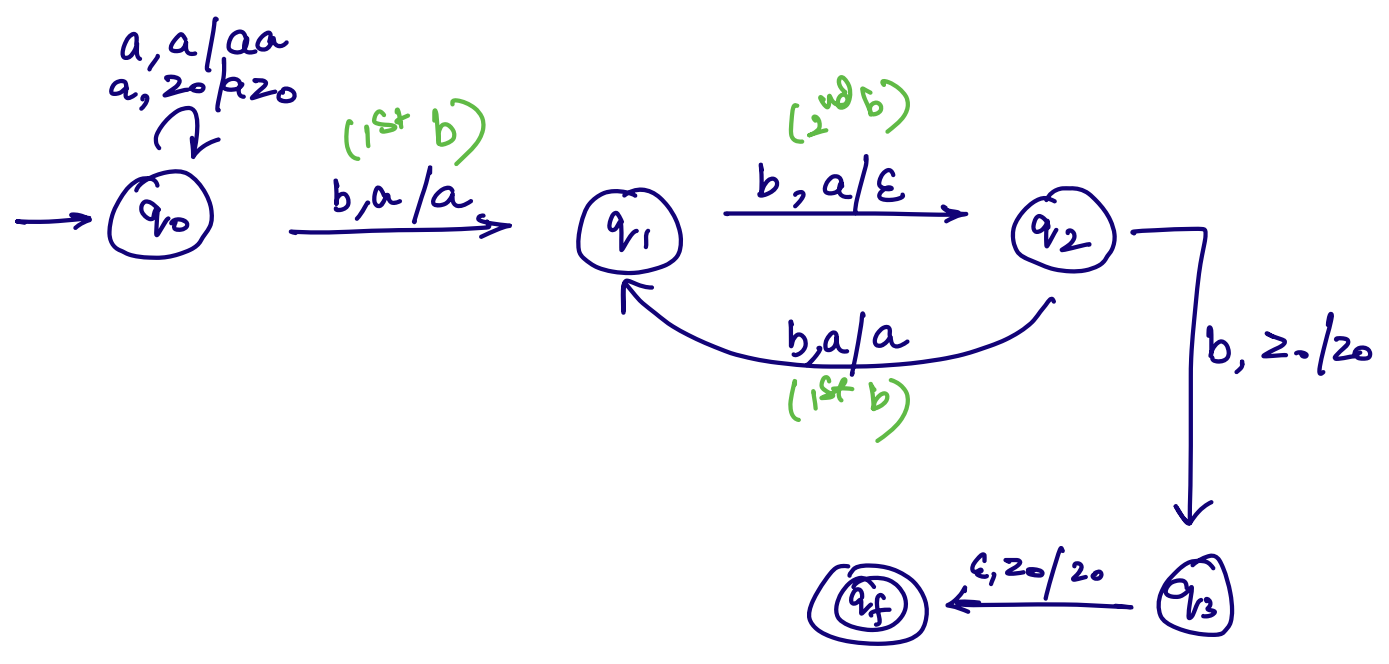
Eg: $a^n b^{2n+1} \quad | \quad n > 1$

$a^n \cdot b^{2n} \cdot b$

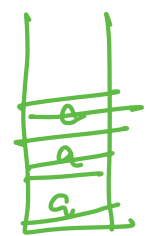
Way 3:



way 2:



eg: $a^n b^n c^n \mid n \geq 1$

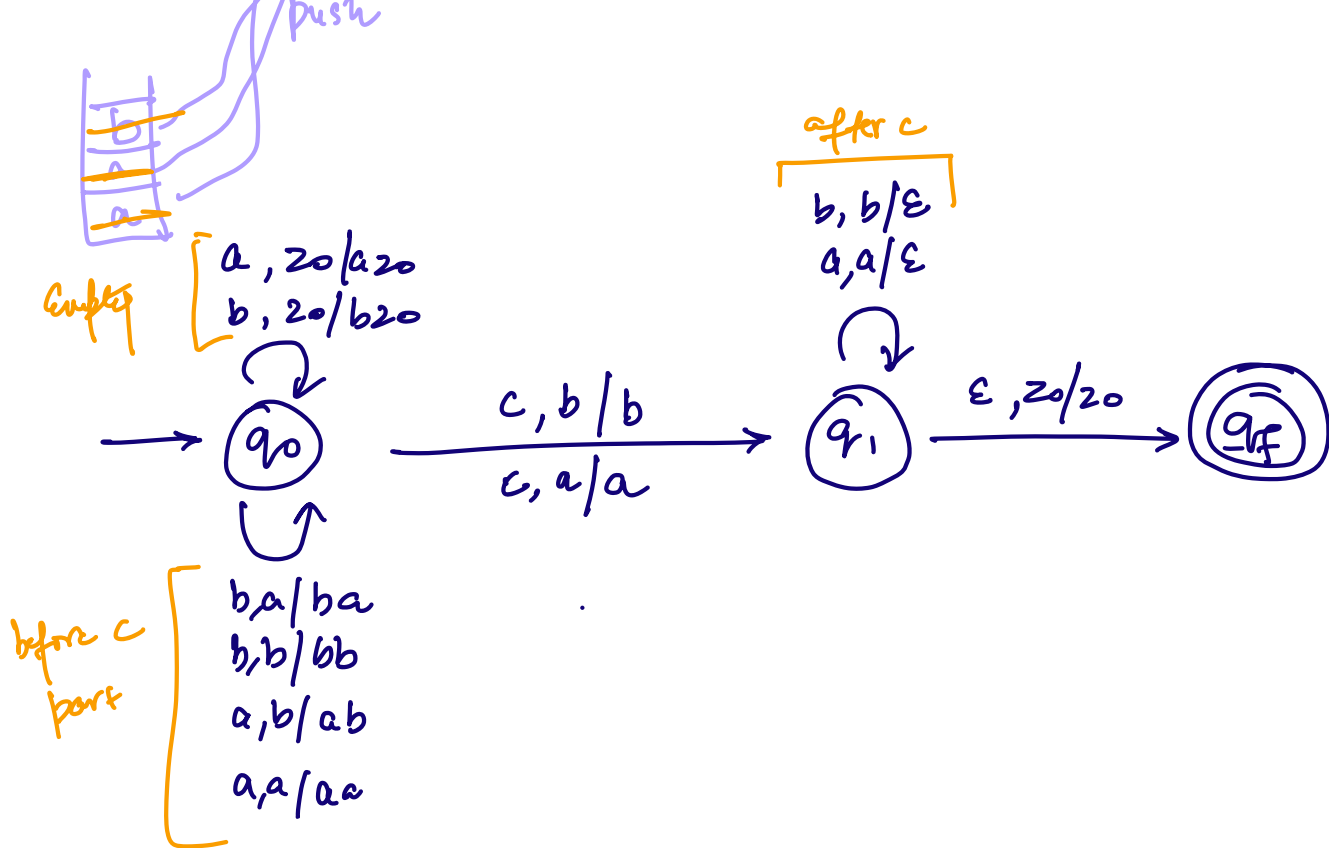


PDA Not possible

eg: $w c w^R \mid w \in (a, b)^+$

Set of all palindromes of odd length





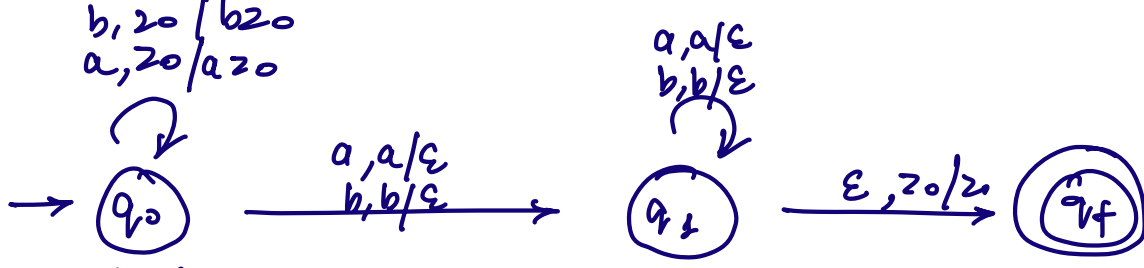
Eg: $ww^R \mid w \in (a,b)^+$
 Even length palindromic strings

Push Pop
abbbba abbbba
w w^R

1. abb/bba
 ← →

2. $a \mid bbbba$
 Non deterministic PDA

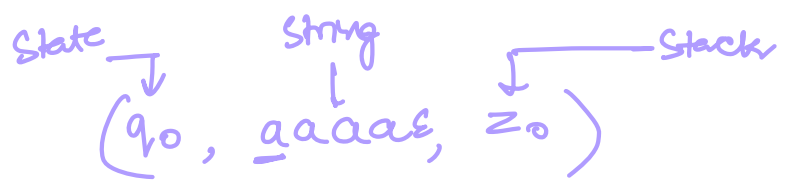
Center push pop
 not a center push push



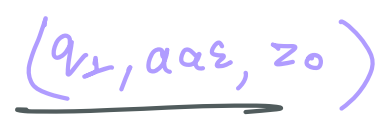
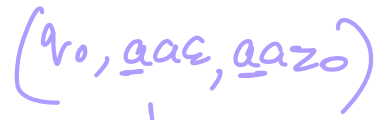
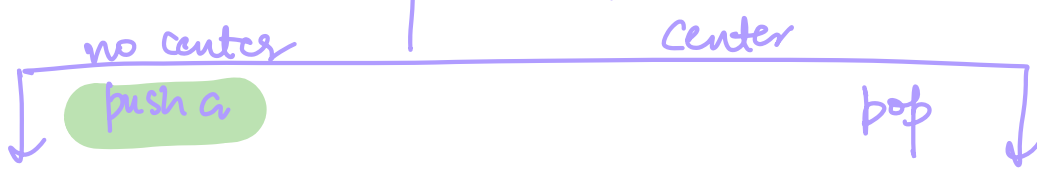
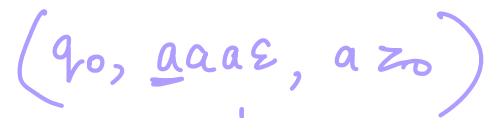
no way
it can be
a center

- b, a/ba
- a, b/ab
- a, a/aa
- b, b/bb

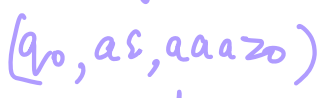
not a
center



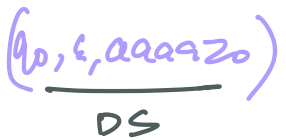
push a



Stack D.S.



pop



DS

DS

